

## 13123 Christmas Lights

Bill has a Christmas Lights String to decorate his house, made with  $K$  lights  $L[1], L[2], \dots, L[K]$  attached in sequence to a wire. The behavior of each light is determined by a programmable controller connected to the wire, turning *on* and *off* lights at every second.

Bill has programmed the controller to change the state of the lights during  $M$  seconds. He defines a pair of numbers  $a_i, b_i$  with  $a_i \leq b_i$ , for each second  $i$  ( $1 \leq i \leq M$ ). At second 0, the string of lights is initialized with a random *initial configuration* (some lights *on* and other lights *off*). At each second  $i$ , from 1 to  $M$ , the state of all lights in  $L[a_i .. b_i]$  is simultaneously switched from *on* to *off* and vice versa. However, Bill added a curious little feature to the controller's algorithm: whenever the ends  $L[a_i]$  or  $L[b_i]$  are *off*, just before the above-described switching takes place at time  $i$ , some more lights in the string can switch states at moment  $i$ . In particular, if  $L[a_i]$  is *off* and there is a light, say at  $l_i$ , to the left of  $a_i$  that is *on* (and all the lights between  $l_i$  and  $a_i$  are *off*), then the lights in the interval  $L[l_i .. a_i - 1]$  will also switch states at moment  $i$ . Similarly, if  $L[b_i]$  is *off* and there is a light, say at  $r_i$ , to the right of  $b_i$  that is *on* (and all the lights between  $b_i$  and  $r_i$  are *off*), then the lights in the interval  $L[b_i + 1 .. r_i]$  will also switch states at moment  $i$ .

Suppose that a light turned *on* is represented with '1' and a light turned *off* is represented with '0'. For example, consider  $K = 18$ ,  $M = 5$ ,  $a_1 = 5$ ,  $b_1 = 12$ ,  $a_2 = 10$ ,  $b_2 = 11$ ,  $a_3 = 5$ ,  $b_3 = 8$ ,  $a_4 = 3$ ,  $b_4 = 6$ ,  $a_5 = 1$ , and  $b_5 = 17$ , with initial configuration 000110010011100000. Note that the state of all lights at each second is:

- 000110010011100000 at second 0.
- 000101101100100000 at second 1.
- 000101101011000000 at second 2.
- 000010010011000000 at second 3.
- 001101100011000000 at second 4.
- 110010011100111110 at second 5.

After several days of operation, Bill suspects that he has created a truly awesome algorithm. For this purpose, he would like to run multiple trials, with different initial configurations and parameters  $a, b$ , but he is afraid the lights will break due to heavy abuse. Can you help him in building an algorithm for finding the final state of all lights at second  $M$  after each trial?

### Input

The first line of the input contains a positive integer  $T$  indicating the number of test cases. The first line of a test case contains two blank-separated integers  $K$  and  $M$  ( $2 \leq K \leq 10^6$ ,  $0 \leq M \leq 10^4$ ) indicating, respectively, the number of lights in the string and the number of seconds to consider. The second line contains a hexadecimal string (using digits '0123456789ABCDEF') without leading zeros, describing the initial configuration of lights if it is written in binary notation. If the given hexadecimal string requires less than  $K$  bits in binary notation, then complete it with leading zeros to reach  $K$  digits. Finally follow  $M$  lines: line  $i$  contains exactly two blank-separated integers  $a_i$  and  $b_i$  describing the parameters controlling the behavior of the lights at second  $i$  ( $1 \leq i \leq M$ ,  $1 \leq a_i \leq b_i \leq K$ ).

## Output

For each test case, print a single line with a hexadecimal string (using digits '0123456789ABCDEF') without leading zeros, describing the state of all lights at second  $M$ . You must use the same notation used to codify the initial configuration of lights.

## Sample Input

```
3
18 5
64E0
5 12
10 11
5 8
3 6
1 17
18 0
0
18 1
0
13 16
```

## Sample Output

```
3273E
0
3C
```